

Computer Science

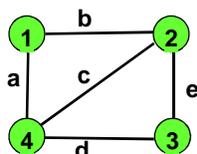
- Lecture notes
- CSAUTFALL2010
- Use by ONLY Group Members

Wintert, 2010

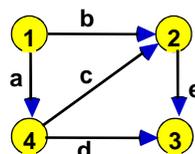
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Notation and Terminology

Network terminology as used in AMO.



An Undirected Graph or
Undirected Network



A Directed Graph or
Directed Network

Network $G = (N, A)$

Node set $N = \{1, 2, 3, 4\}$

Arc Set $A = \{(1,2), (1,3), (3,2), (3,4), (2,4)\}$

In an undirected graph, $(i,j) = (j,i)$

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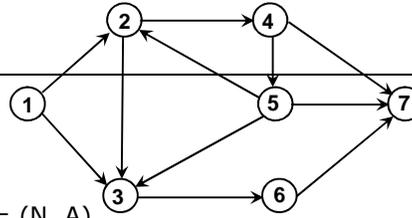
<p>Path: Example: 5, 2, 3, 4. (or 5, c, 2, b, 3, e, 4) No node is repeated. Directions are ignored.</p>	
<p>Directed Path: Example: 1, 2, 5, 3, 4 (or 1, a, 2, c, 5, d, 3, e, 4) No node is repeated. Directions are important.</p>	
<p>Cycle (or circuit or loop): 1, 2, 3, 1. (or 1, a, 2, b, 3, e) A path with 2 or more nodes, except that the first node is the last node. Directions are ignored.</p>	
<p>Directed Cycle: (1, 2, 3, 4, 1) or 1, a, 2, b, 3, c, 4, d, 1 No node is repeated. Directions are important.</p>	

Walks

Walks are paths that can repeat nodes and arcs
 Example of a **directed walk**: 1-2-3-5-4-2-3-5
 A walk is **closed** if its first and last nodes are the same.
 A closed walk is a cycle except that it can repeat nodes and arcs.

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Network Notation



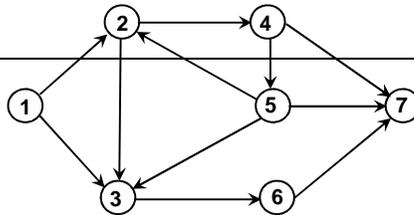
Network (or graph) $G = (N, A)$

N : Node Set = $\{1, 2, 3, 4, 5, 6, 7\}$

A : Arc Set = $\{(1, 2), (1, 3), (2, 3), (2, 4), (3, 6), (4, 5), (4, 7), (5, 2), (5, 7), (6, 7)\}$

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Tails and Head Nodes

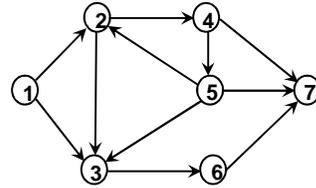


- Tail Node : Node i is the tail node of arc (i, j)
- Head Node : Node j is the head node of arc (i, j)
- Endpoints : Nodes i and j are endpoints of arc (i, j)
- Arc Adjacency List $A(i) = \{(i, j) : (i, j) \in A\}$
- Observe that $\sum_{i \in N} |A(i)| = m$

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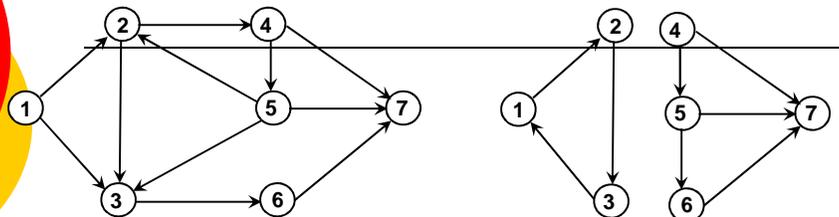
Subgraph

- Subgraph: $G' = (N', A')$ is a subgraph of $G = (N, A)$ if $N' \subseteq N$ and $A' \subseteq A$.
- Spanning Subgraph: $G' = (N', A')$ is a spanning subgraph of $G = (N, A)$ if $N' = N$ and $A' \subseteq A$.
- Acyclic Network: A network is acyclic if it contains no directed cycle



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Network Connectivity

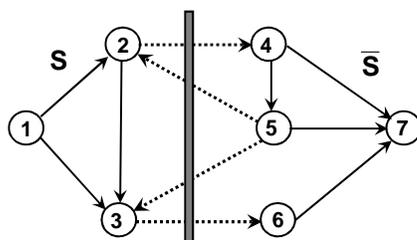


- Connected Graph: If every pair of nodes is connected by a path.
- Strongly Connected Graph: If every pair of nodes is connected by a *directed* path.
- Disconnected Graph: If some pair of nodes is not connected by a path. Maximally connected subgraphs of a disconnected graph are called components.

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Cuts

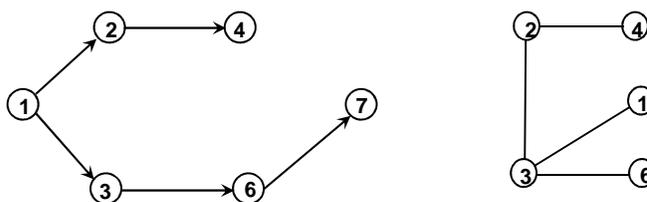
- Cut: A cut is a set of arcs whose deletion disconnects the network into two components S and \bar{S} and no subset of it has this property. A cut is represented by $[S, \bar{S}]$.



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Trees

- ~~Tree: A tree is a connected graph that contains no cycle.~~

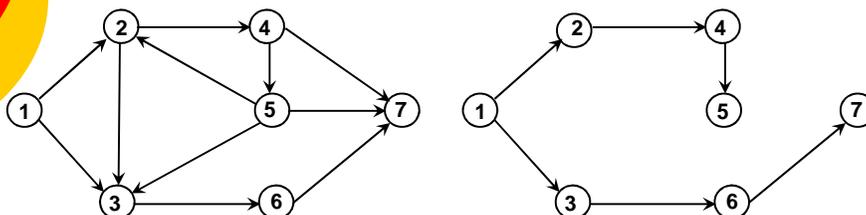


- A tree on n nodes contains exactly $(n-1)$ arcs.
- A tree has at least two leaf nodes.
- Every two nodes of a tree are connected by a unique path.

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Spanning Trees

- Spanning Tree: A subgraph of a graph which is a tree and spans all nodes of the graph.

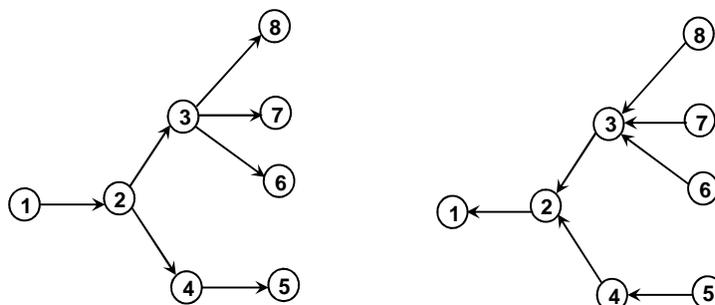


- Forest: A graph that contains no cycle is a forest.
- A forest on n nodes and k components has $n-k$ arcs.

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Directed Out-Trees and In-Trees

- Directed-out-tree: A tree in which the unique path in the tree from node s to every other node is a directed path.
- Directed-in-tree: A tree in which the unique path in the tree from any node to node s is a directed path.

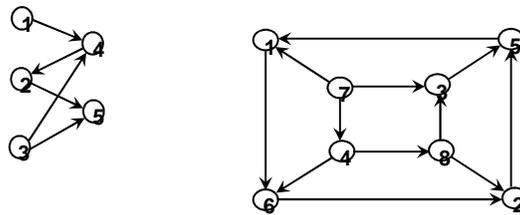


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Bipartite Graph

- Bipartite Graph: A graph $G = (N, A)$ is a bipartite graph if node set N can be partitioned into the sets N^1 and N^2 such that each arc (i, j) has either

(i) $i \in N^1$ and $j \in N^2$; or (ii) $i \in N^2$ and $j \in N^1$.



- A graph G is a bipartite graph if and only if every cycle in it has even number of arcs.

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Network Representations

- Role of data structures is crucial in developing efficient algorithms. The manner in which networks are stored in a computer determines the performance of a network algorithm.
- We need to perform some of the following operations:
 - Scan through the node list
 - Scan through the arc list
 - Determine all arcs emanating from a specific node i
 - Determine all arcs entering a specific node i
 - Determine all arcs emanating from or entering a specific node i

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Common Network Representations

- Node-arc incidence matrix
- Node-node adjacency matrix
- Adjacency list
- Forward star representation
- Reverse star representation
- Forward and reverse star representations

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Network representations

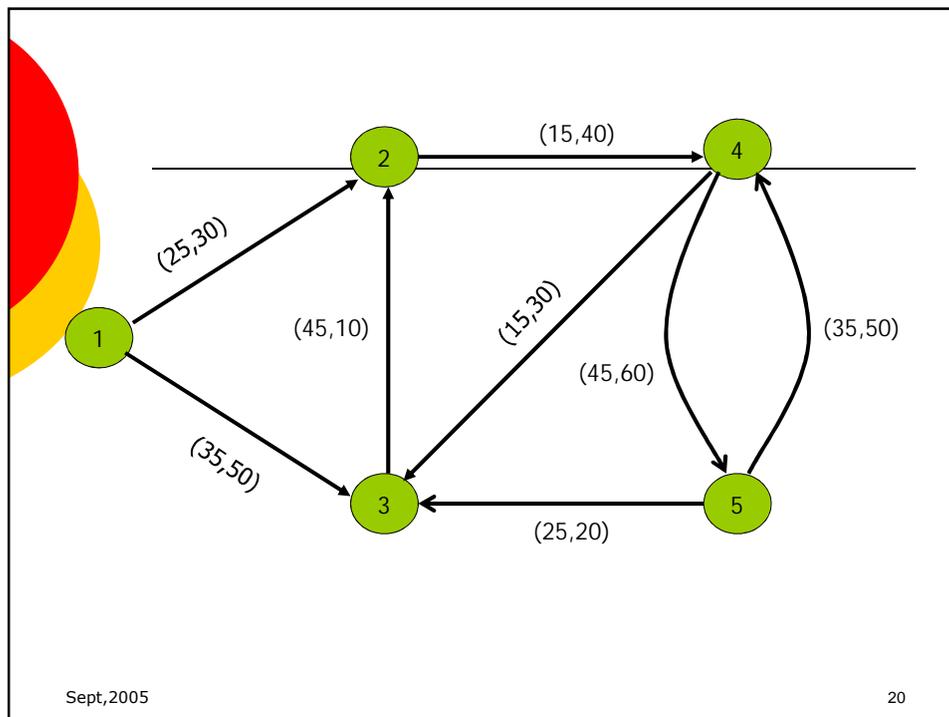
- We typically need to store two types of information:
 - Network topology
 - Nodes and arc structures
 - Data
 - Costs
 - Capacities
 - Supplies / demands

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Incidence matrix representation

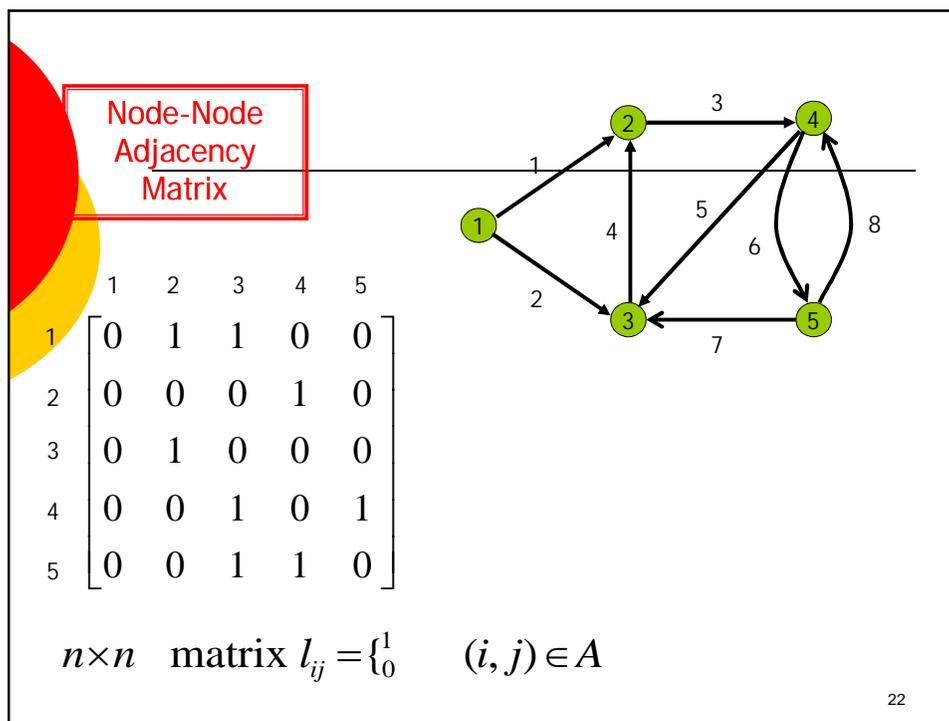
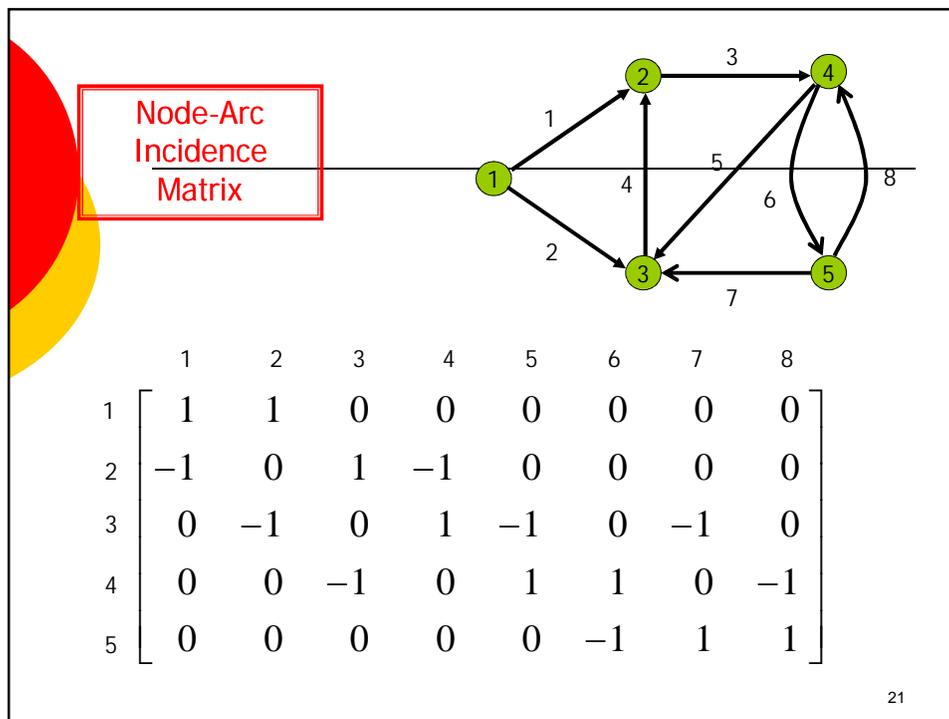
- N $n \times m$ matrix contains
 - One row for each node of the network
 - One column for each arc
 - Column (i, j) has only two nonzero elements
 - +1 in row i
 - -1 in row j

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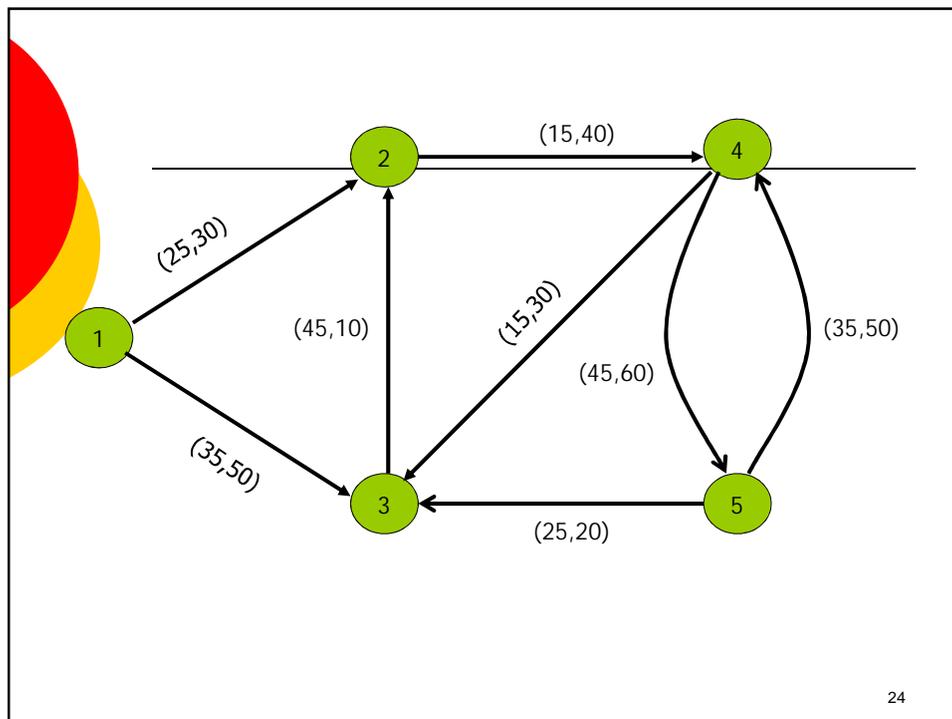
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Adjacency list

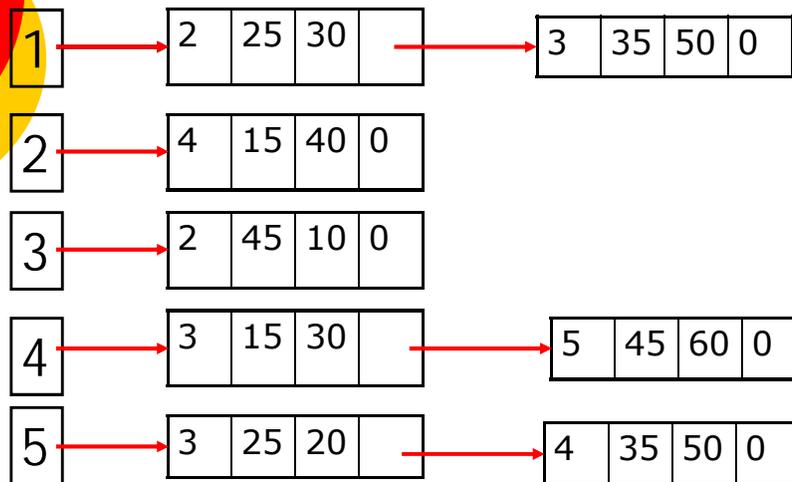
- The adjacency list representation stores the node adjacency list of each node as a singly linked list.

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Adjacency Lists



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Forward Star Representation

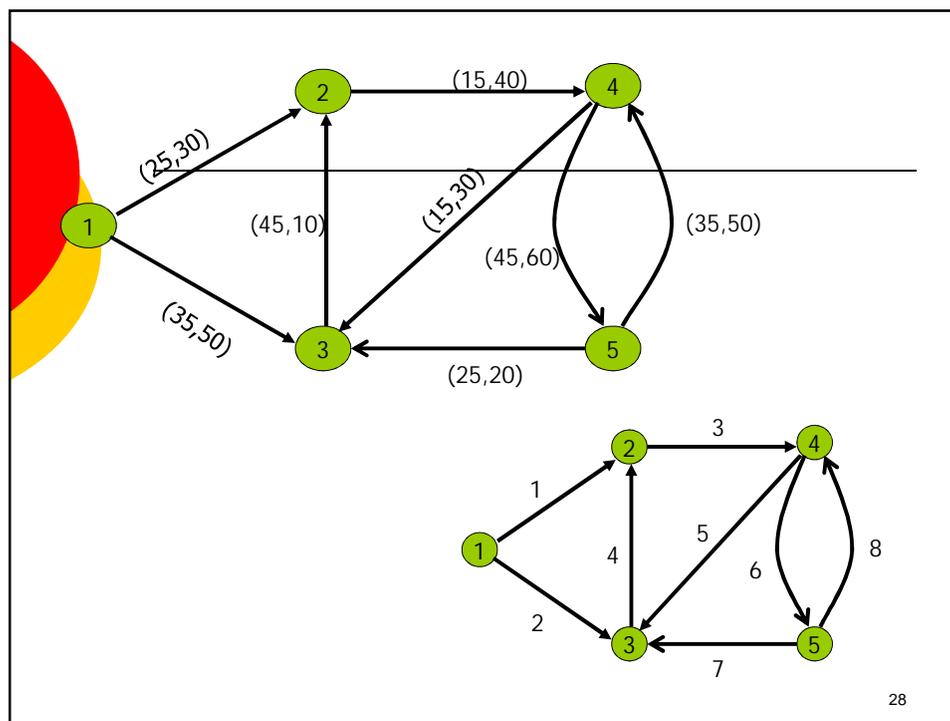
- Similar to the adjacency list representation which stores the node adjacency list of each node
- It stores them in a single array
- We number the arcs in a specific order
- We store the tails, heads, costs, and capacities of the arcs in four arrays: tail, head, cost, and capacity

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Forward Star Representation

- We maintain a pointer with each node i , denoted by $\text{Point}(i)$,
- $\text{Point}(i) :=$ the smallest number arc in the arc list that emanates from node i
- $\text{Point}(i) := \text{Point}(i+1)$
 - if node i has no outgoing arcs
- $\text{Point}(1) = 1$
- $\text{Point}(n+1) = m+1$

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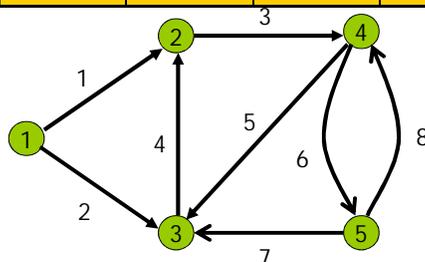
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Forward Star Representation

	point		tail	head	cost	capacity
1	1	1	1	2	25	30
2	3	2	1	3	35	50
3	4	3	2	4	15	40
4	5	4	3	2	45	10
5	7	5	4	3	15	30
6	7	6	4	5	45	60
7	9	7	5	3	25	20
8	9	8	5	4	35	50

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	Point		tail	head	cost	capacity
1	1	1	1	2	25	30
2	3	2	1	3	35	50
3	4	3	2	4	15	40
4	5	4	3	2	45	10
5	7	5	4	3	15	30
6	7	6	4	5	45	60
7	9	7	5	3	25	20
8	9	8	5	4	35	50



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○ Note: the FSR will store the outgoing arcs of node i at positions

- Point (i) to point $(i+1) - 1$
- In the arc list
- If $\text{point}(i) > \text{point}(i+1)-1$ node i has no outgoing arc
- FSR provides us with an efficient way for determining the set of outgoing arcs of any node

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Exercises due date Sunday 5/7/89

- Exercises
- 2.4, 2.8, 2.10
- 2.12 , 2.13, 2.16
- 2.19, 2.20, 2.23 , 2.24
- 2.32, 2.33, 2.34, 2.35 2.36 , 2.37, 2.38

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